### Queue:

A queue is another special kind of list, where items are inserted at one end called the rear and deleted at the other end called the front. Another name for a queue is a “FIFO” or “First-in-first-out” list.

The operations for a queue are analogues to those for a stack, the difference is that the insertions go at the end of the list, rather than the beginning. We shall use the following operations on queues:

* + - *enqueue*: which inserts an element at the end of the queue.
    - *dequeue*: which deletes an element at the start of the queue.

### Representation of Queue:

Let us consider a queue, which can hold maximum of five elements. Initially the queue is empty.

0 1 2 3 4

|  |  |  |  |  |
| --- | --- | --- | --- | --- |
|  |  |  |  |  |

F R

Q u eu e E mp t y

F RO NT = RE A R = 0

Now, insert 11 to the queue. Then queue status will be:

0 1 2 3 4

|  |  |  |  |  |
| --- | --- | --- | --- | --- |
| 11 |  |  |  |  |

F R

RE A R = RE A R + 1 = 1 FRO NT = 0

Next, insert 22 to the queue. Then the queue status is:

0 1 2 3 4

|  |  |  |  |  |
| --- | --- | --- | --- | --- |
| 11 | 22 |  |  |  |

F R

RE A R = RE A R + 1 = 2 FRO NT = 0

Again insert another element 33 to the queue. The status of the queue is:

0 1 2 3 4

|  |  |  |  |  |
| --- | --- | --- | --- | --- |
| 11 | 22 | 33 |  |  |

F R

RE A R = RE A R + 1 = 3 FRO NT = 0

Now, delete an element. The element deleted is the element at the front of the queue. So the status of the queue is:

0 1 2 3 4

|  |  |  |  |  |
| --- | --- | --- | --- | --- |
|  | 22 | 33 |  |  |

RE A R = 3

F RO NT = F R O NT + 1 = 1

F R

Again, delete an element. The element to be deleted is always pointed to by the FRONT pointer. So, 22 is deleted. The queue status is as follows:

0 1 2 3 4

|  |  |  |  |  |
| --- | --- | --- | --- | --- |
|  |  | 33 |  |  |

F R

RE A R = 3

F RO NT = F R O NT + 1 = 2

Now, insert new elements 44 and 55 into the queue. The queue status is:

0 1 2 3 4

|  |  |  |  |  |
| --- | --- | --- | --- | --- |
|  |  | 33 | 44 | 55 |

F R

RE A R = 5 FRO NT = 2

Next insert another element, say 66 to the queue. We cannot insert 66 to the queue as the rear crossed the maximum size of the queue (i.e., 5). There will be queue full signal. The queue status is as follows:

0 1 2 3 4

|  |  |  |  |  |
| --- | --- | --- | --- | --- |
|  |  | 33 | 44 | 55 |

RE A R = 5 FRO NT = 2

F R

Now it is not possible to insert an element 66 even though there are two vacant positions in the linear queue. To over come this problem the elements of the queue are to be shifted towards the beginning of the queue so that it creates vacant position at the rear end. Then the FRONT and REAR are to be adjusted properly. The element 66 can be inserted at the rear end. After this operation, the queue status is as follows:

0 1 2 3 4

|  |  |  |  |  |
| --- | --- | --- | --- | --- |
| 33 | 44 | 55 | 66 |  |

F R

RE A R = 4 FRO NT = 0

This difficulty can overcome if we treat queue position with index 0 as a position that comes after position with index 4 i.e., we treat the queue as a **circular queue.**

### Source code for Queue operations using array:

In order to create a queue we require a one dimensional array Q(1:n) and two variables *front* and *rear*. The conventions we shall adopt for these two variables are that *front* is always 1 less than the actual front of the queue and rear always points to the last element in the queue. Thus, front = rear if and only if there are no elements in the queue. The initial condition then is front = rear = 0. The various queue operations to perform creation, deletion and display the elements in a queue are as follows:

* + - 1. insertQ(): inserts an element at the end of queue Q.
      2. deleteQ(): deletes the first element of Q.
      3. displayQ(): displays the elements in the queue.

# include <conio.h> # define MAX 6

int Q[MAX]; int front, rear;

void insertQ()

{

int data;

if(rear == MAX)

{

}

else

{

}

}

printf("\n Linear Queue is full"); return;

printf("\n Enter data: "); scanf("%d", &data); Q[rear] = data;

rear++;

printf("\n Data Inserted in the Queue ");

void deleteQ()

{

if(rear == front)

{

}

else

{

}

}

printf("\n\n Queue is Empty.."); return;

printf("\n Deleted element from Queue is %d", Q[front]); front++;

void displayQ()

{

int i;

if(front == rear)

{

}

else

{

printf("\n\n\t Queue is Empty"); return;

printf("\n Elements in Queue are: "); for(i = front; i < rear; i++)

{

printf("%d\t", Q[i]);

}

}

}

int menu()

{

int ch; clrscr();

printf("\n \tQueue operations using ARRAY.."); printf("\n -----------\*\*\*\*\*\*\*\*\*\* \n");

printf("\n 1. Insert "); printf("\n 2. Delete "); printf("\n 3. Display"); printf("\n 4. Quit ");

printf("\n Enter your choice: "); scanf("%d", &ch);

return ch;

}

void main()

{

int ch; do

{

ch = menu(); switch(ch)

{

case 1:

case 2:

case 3:

case 4:

}

insertQ(); break;

deleteQ(); break;

displayQ(); break;

return;

getch();

} while(1);

}

### Linked List Implementation of Queue:

We can represent a queue as a linked list. In a queue data is deleted from the front end and inserted at the rear end. We can perform similar operations on the two ends of a list. We use two pointers *front* and *rear* for our linked queue implementation.

The linked queue looks as shown in figure 4.4:

**front**

100

**rear**

400

10 200

100

20 300

200

30 400 40 **X**

300 400

Figure 4.4. Linked Queue representation

### Source code for queue operations using linked list:

# include <stdlib.h> # include <conio.h>

struct queue

{

int data;

struct queue \*next;

};

typedef struct queue node; node \*front = NULL;

node \*rear = NULL;

node\* getnode()

{

node \*temp;

temp = (node \*) malloc(sizeof(node)) ; printf("\n Enter data ");

scanf("%d", &temp -> data); temp -> next = NULL;

return temp;

}

void insertQ()

{

node \*newnode; newnode = getnode(); if(newnode == NULL)

{

printf("\n Queue Full"); return;

}

if(front == NULL)

{

}

else

{

}

front = newnode; rear = newnode;

rear -> next = newnode; rear = newnode;

printf("\n\n\t Data Inserted into the Queue..");

}

void deleteQ()

{

node \*temp; if(front == NULL)

{

printf("\n\n\t Empty Queue.."); return;

}

temp = front;

front = front -> next;

printf("\n\n\t Deleted element from queue is %d ", temp -> data); free(temp);

}

void displayQ()

{

node \*temp; if(front == NULL)

{

}

else

{

printf("\n\n\t\t Empty Queue ");

temp = front;

printf("\n\n\n\t\t Elements in the Queue are: "); while(temp != NULL )

{

printf("%5d ", temp -> data); temp = temp -> next;

}

}

}

char menu()

{

char ch; clrscr();

printf("\n \t..Queue operations using pointers.. "); printf("\n\t -----------\*\*\*\*\*\*\*\*\*\* \n");

printf("\n 1. Insert "); printf("\n 2. Delete "); printf("\n 3. Display"); printf("\n 4. Quit ");

printf("\n Enter your choice: "); ch = getche();

return ch;

}

void main()

{

char ch; do

{

ch = menu(); switch(ch)

{

case '1' :

insertQ(); break;

case '2' :

deleteQ(); break;

case '3' :

displayQ(); break;

case '4':

return;

}

getch();

} while(ch != '4');

}

### Applications of Queue:

1. It is used to schedule the jobs to be processed by the CPU.
2. When multiple users send print jobs to a printer, each printing job is kept in the printing queue. Then the printer prints those jobs according to first in first out (FIFO) basis.
3. Breadth first search uses a queue data structure to find an element from a graph.

### Circular Queue:

A more efficient queue representation is obtained by regarding the array Q[MAX] as circular. Any number of items could be placed on the queue. This implementation of a queue is called a circular queue because it uses its storage array as if it were a circle instead of a linear list.

There are two problems associated with linear queue. They are:

* Time consuming: linear time to be spent in shifting the elements to the beginning of the queue.
* Signaling queue full: even if the queue is having vacant position.

For example, let us consider a linear queue status as follows:

0 1 2 3 4

|  |  |  |  |  |
| --- | --- | --- | --- | --- |
|  |  | 33 | 44 | 55 |

F R

RE A R = 5 FRO NT = 2

Next insert another element, say 66 to the queue. We cannot insert 66 to the queue as the rear crossed the maximum size of the queue (i.e., 5). There will be queue full signal. The queue status is as follows:

0 1 2 3 4

|  |  |  |  |  |
| --- | --- | --- | --- | --- |
|  |  | 33 | 44 | 55 |

F R

RE A R = 5 FRO NT = 2

This difficulty can be overcome if we treat queue position with index zero as a position that comes after position with index four then we treat the queue as a **circular queue.**

In circular queue if we reach the end for inserting elements to it, it is possible to insert new elements if the slots at the beginning of the circular queue are empty.

### Representation of Circular Queue:

Let us consider a circular queue, which can hold maximum (MAX) of six elements. Initially the queue is empty.

F R

1 Q u eu e E mp t y

5

0

3

2

4 MA X = 6

F RO NT = RE A R = 0 CO U NT = 0

Circ u lar Q u e u e

Now, insert 11 to the circular queue. Then circular queue status will be:

F

R

5

0

11

1

3

2

FRO NT = 0

4 RE A R = ( RE A R + 1) % 6 = 1

CO U NT = 1

Circ u lar Q u e u e

Insert new elements 22, 33, 44 and 55 into the circular queue. The circular queue status is:

4 1 FRONT = 0

F

R

5

0

11

55

22

44

33

3

2

REAR = (REAR + 1) % 6 = 5

COUNT = 5

Circular Queue

Now, delete an element. The element deleted is the element at the front of the circular queue. So, 11 is deleted. The circular queue status is as follows:

R

F

5

0

55

22 1

44

33

3

2

F RO NT = (F R O NT + 1) % 6 = 1

4 RE A R = 5

CO U NT = CO U NT - 1 = 4

Circ u lar Q u e u e

Again, delete an element. The element to be deleted is always pointed to by the FRONT pointer. So, 22 is deleted. The circular queue status is as follows:

R

F RO NT = (F R O NT + 1) % 6 = 2

5

0

55

1

44

33

3

2

4 RE A R = 5

CO U NT = CO U NT - 1 = 3

F

Circ u lar Q u e u e

Again, insert another element 66 to the circular queue. The status of the circular queue is:

R

4 FRO NT = 2

5

0

66

55

1

44

33

3

2

RE A R = ( RE A R + 1) % 6 = 0 C O U NT = C O U NT + 1 = 4

F

Circ u lar Q u e u e

Now, insert new elements 77 and 88 into the circular queue. The circular queue status is:

5 0

66 77

4 55

3

88 1

44 33

2 R

F

F RO NT = 2, RE A R = 2 RE A R = RE A R % 6 = 2 CO U NT = 6

Circ u lar Q u e u e

Now, if we insert an element to the circular queue, as COUNT = MAX we cannot add the element to circular queue. So, the circular queue is *full*.

### Source code for Circular Queue operations, using array:

# include <stdio.h> # include <conio.h> # define MAX 6

int CQ[MAX];

int front = 0; int rear = 0; int count = 0;

void insertCQ()

{

int data;

if(count == MAX)

{

}

else

{

}

}

printf("\n Circular Queue is Full");

printf("\n Enter data: "); scanf("%d", &data); CQ[rear] = data;

rear = (rear + 1) % MAX; count ++;

printf("\n Data Inserted in the Circular Queue ");

void deleteCQ()

{

if(count == 0)

{

}

else

{

}

}

printf("\n\nCircular Queue is Empty..");

printf("\n Deleted element from Circular Queue is %d ", CQ[front]); front = (front + 1) % MAX;

count --;

void displayCQ()

{

int i, j; if(count == 0)

{

}

else

{

printf("\n\n\t Circular Queue is Empty ");

printf("\n Elements in Circular Queue are: "); j = count;

for(i = front; j != 0; j--)

{

printf("%d\t", CQ[i]); i = (i + 1) % MAX;

}

}

}

int menu()

{

int ch; clrscr();

printf("\n \t Circular Queue Operations using ARRAY.."); printf("\n -----------\*\*\*\*\*\*\*\*\*\* \n");

printf("\n 1. Insert "); printf("\n 2. Delete "); printf("\n 3. Display"); printf("\n 4. Quit ");

printf("\n Enter Your Choice: "); scanf("%d", &ch);

return ch;

}

void main()

{

int ch; do

{

ch = menu(); switch(ch)

{

case 1:

case 2:

case 3:

case 4: default:

}

insertCQ(); break;

deleteCQ(); break;

displayCQ(); break;

return;

printf("\n Invalid Choice ");

getch();

} while(1);

}

### Deque:

In the preceding section we saw that a queue in which we insert items at one end and from which we remove items at the other end. In this section we examine an extension of the queue, which provides a means to insert and remove items at both ends of the queue. This data structure is a *deque*. The word *deque* is an acronym derived from *double-ended queue*. Figure 4.5 shows the representation of a deque.

Deletion

Insertion

Insertion

Deletion

front rear

|  |  |  |  |  |
| --- | --- | --- | --- | --- |
| 36 | 16 | 56 | 62 | 19 |

Figure 4.5. Representation of a deque.

A deque provides four operations. Figure 4.6 shows the basic operations on a deque.

* + - enqueue\_front: insert an element at front.
    - dequeue\_front: delete an element at front.
    - enqueue\_rear: insert element at rear.
    - dequeue\_rear: delete element at rear.

|  |  |  |  |
| --- | --- | --- | --- |
|  | 11 | 22 |  |

|  |  |  |  |  |
| --- | --- | --- | --- | --- |
| ) | 33 | 11 | 22 |  |

|  |  |  |  |  |
| --- | --- | --- | --- | --- |
|  | 33 | 11 | 22 | 44 |

|  |  |  |  |  |
| --- | --- | --- | --- | --- |
| 5) |  | 11 | 22 | de |

|  |  |  |  |  |  |
| --- | --- | --- | --- | --- | --- |
|  |  | 11 | 22 | 44 |  |

Figure 4.6. Basic operations on deque

enqueue\_front(33

enqueue\_rear(44)

dequeue\_front(33)

enqueue\_front(5

queue\_rear(44)

22

11

55

There are two variations of deque. They are:

* + - Input restricted deque (IRD)
    - Output restricted deque (ORD)

An Input restricted deque is a deque, which allows insertions at one end but allows deletions at both ends of the list.

An output restricted deque is a deque, which allows deletions at one end but allows insertions at both ends of the list.

### Priority Queue:

A priority queue is a collection of elements such that each element has been assigned a priority and such that the order in which elements are deleted and processed comes from the following rules:

1. An element of higher priority is processed before any element of lower priority.
2. two elements with same priority are processed according to the order in which they were added to the queue.

A prototype of a priority queue is time sharing system: programs of high priority are processed first, and programs with the same priority form a standard queue. An efficient implementation for the Priority Queue is to use heap, which in turn can be used for sorting purpose called heap sort.

### Exercises

1. What is a linear data structure? Give two examples of linear data structures.
2. Is it possible to have two designs for the same data structure that provide the same functionality but are implemented differently?
3. What is the difference between the logical representation of a data structure and the physical representation?
4. Transform the following infix expressions to reverse polish notation:
   1. A  B \* C – D + E / F / (G + H)

b) ((A + B) \* C – (D – E))  (F + G)

c) A – B / (C \* D  E)

d) (a + b  c  d) \* (e + f / d)) f) 3 – 6 \* 7 + 2 / 4 \* 5 – 8

g) (A – B) / ((D + E) \* F)

h) ((A + B) / D)  ((E – F) \* G)

1. Evaluate the following postfix expressions: a) P1: 5, 3, +, 2, \*, 6, 9, 7, -, /, -

b) P2: 3, 5, +, 6, 4, -, \*, 4, 1, -, 2, , +

c) P3 : 3, 1, +, 2, , 7, 4, -, 2, \*, +, 5, -

1. Consider the usual algorithm to convert an infix expression to a postfix expression. Suppose that you have read 10 input characters during a conversion and that the stack now contains these symbols:

bottom

+ (

\*

Now, suppose that you read and process the 11th symbol of the input. Draw the stack for the case where the 11th symbol is:

1. A number:
2. A left parenthesis:
3. A right parenthesis:
4. A minus sign:
5. A division sign:
6. Write a program using stack for parenthesis matching. Explain what modifications would be needed to make the parenthesis matching algorithm check expressions with different kinds of parentheses such as (), [] and {}'s.
7. Evaluate the following prefix expressions: a) + \* 2 + / 14 2 5 1

b) - \* 6 3 – 4 1

c) + + 2 6 + - 13 2 4

1. Convert the following infix expressions to prefix notation: a) ((A + 2) \* (B + 4)) -1

b) Z – ((((X + 1) \* 2) – 5) / Y)

c) ((C \* 2) + 1) / (A + B)

d) ((A + B) \* C – (D - E))  (F + G)

1. A – B / (C \* D  E)
2. Write a “C” function to copy one stack to another assuming
   1. The stack is implemented using array.
   2. The stack is implemented using linked list.
3. Write an algorithm to construct a fully parenthesized infix expression from its postfix equivalent. Write a “C” function for your algorithm.
4. How can one convert a postfix expression to its prefix equivalent and vice-versa?
5. A double-ended queue (deque) is a linear list where additions and deletions can be performed at either end. Represent a deque using an array to store the elements of the list and write the “C” functions for additions and deletions.
6. In a circular queue represented by an array, how can one specify the number of elements in the queue in terms of “front”, “rear” and MAX-QUEUE-SIZE? Write a “C” function to delete the K-th element from the “front” of a circular queue.
7. Can a queue be represented by a circular linked list with only one pointer pointing to the tail of the queue? Write “C” functions for the “add” and “delete” operations on such a queue
8. Write a “C” function to test whether a string of opening and closing parenthesis is well formed or not.
9. Represent N queues in a single one-dimensional array. Write functions for “add” and “delete” operations on the ith queue
10. Represent a stack and queue in a single one-dimensional array. Write functions for “push”, “pop” operations on the stack and “add”, “delete” functions on the queue.

### Multiple Choice Questions

1. Which among the following is a linear data structure: [ D ]
   1. Queue
   2. Stack
   3. Linked List
   4. all the above
2. Which among the following is a Dynamic data structure: [ A ]
   1. Double Linked List
   2. Queue
   3. Stack
   4. all the above
3. Stack is referred as: [ A ]
   1. Last in first out list
   2. First in first out list
   3. both A and B
   4. none of the above
4. A stack is a data structure in which all insertions and deletions of entries are made at:

[ A ]

* 1. One end
  2. In the middle
  3. Both the ends
  4. At any position

1. A queue is a data structure in which all insertions and deletions are made respectively at:

[ A ]

* 1. rear and front
  2. front and front
  3. front and rear
  4. rear and rear

1. Transform the following infix expression to postfix form: (A + B) \* (C – D) / E

[ D ]

* 1. A B \* C + D / -
  2. A B C \* C D / - +
  3. A B + C D \* - / E
  4. A B + C D - \* E /

1. Transform the following infix expression to postfix form: A - B / (C \* D)

[ B ]

* 1. A B \* C D - /
  2. A B C D \* / -
  3. / - D C \* B A
  4. - / \* A B C D

1. Evaluate the following prefix expression: \* - + 4 3 5 / + 2 4 3 [ A ]
   1. 4
   2. 8
   3. 1
   4. none of the above
2. Evaluate the following postfix expression: 1 4 18 6 / 3 + + 5 / + [ C ]
   1. 8
   2. 2
   3. 3
   4. none of the above
3. Transform the following infix expression to prefix form: ((C \* 2) + 1) / (A + B)

[ B ]

A. A B + 1 2 C \* + /

B. / + \* C 2 1 + A B

C. / \* + 1 2 C A B +

D. none of the above

1. Transform the following infix expression to prefix form: Z – ((((X + 1) \* 2) – 5) / Y)

[ D ]

A. / - \* + X 1 2 5 Y

B. Y 5 2 1 X + \* - /

C. / \* - + X 1 2 5 Y

D. none of the above

1. Queue is also known as: [ B ]
   1. Last in first out list
   2. First in first out list
   3. both A and B
   4. none of the above
2. One difference between a queue and a stack is: [ C ]
   1. Queues require dynamic memory, but stacks do not.
   2. Stacks require dynamic memory, but queues do not.
   3. Queues use two ends of the structure; stacks use only one.
   4. Stacks use two ends of the structure, queues use only one.
3. If the characters 'D', 'C', 'B', 'A' are placed in a queue (in that order), and then removed one at a time, in what order will they be removed?

[ D ]

* 1. ABCD
  2. ABDC
  3. DCAB
  4. DCBA

1. Suppose we have a circular array implementation of the queue class, with ten items in the queue stored at data[2] through data[11]. The CAPACITY is 42. Where does the push member function place the new entry in the array?

[ D ]

* 1. data[1]
  2. data[2]
  3. data[11]
  4. data[12]

1. Consider the implementation of the queue using a circular array. What goes wrong if we try to keep all the items at the front of a partially-filled array (so that data[0] is always the front).
   1. The constructor would require linear time.
   2. The get\_front function would require linear time.
   3. The insert function would require linear time.
   4. The is\_empty function would require linear time.
2. In the linked list implementation of the queue class, where does the push member function place the new entry on the linked list?
   1. At the head
   2. At the tail
   3. After all other entries that are greater than the new entry.
   4. After all other entries that are smaller than the new entry.
3. In the circular array version of the queue class (with a fixed-sized array), which operations require linear time for their worst-case behavior?

[ B ]

[ A ]

[ ]

* 1. front
  2. push
  3. empty
  4. None of these.

1. In the linked-list version of the queue class, which operations require [ ] linear time for their worst-case behavior?
   1. front
   2. push
   3. empty
   4. None of these operations.
2. To implement the queue with a linked list, keeping track of a front pointer and a rear pointer. Which of these pointers will change during an insertion into a NONEMPTY queue?

[ B ]

* 1. Neither changes
  2. Only front\_ptr changes.
  3. Only rear\_ptr changes.
  4. Both change.

1. To implement the queue with a linked list, keeping track of a front pointer and a rear pointer. Which of these pointers will change during an insertion into an EMPTY queue?

[ D ]

* 1. Neither changes
  2. Only front\_ptr changes.
  3. Only rear\_ptr changes.
  4. Both change.

1. Suppose top is called on a priority queue that has exactly two entries with equal priority. How is the return value of top selected?
   1. The implementation gets to choose either one.
   2. The one which was inserted first.
   3. The one which was inserted most recently.
   4. This can never happen (violates the precondition)

[ B ]

1. Entries in a stack are "ordered". What is the meaning of this statement? [ D ]
   1. A collection of stacks can be sorted.
   2. Stack entries may be compared with the '<' operation.
   3. The entries must be stored in a linked list.
   4. There is a first entry, a second entry, and so on.
2. The operation for adding an entry to a stack is traditionally called: [ D ]
   1. add
   2. append
   3. insert
   4. push

|  |  |  |  |  |
| --- | --- | --- | --- | --- |
| 1. The operation for removing an entry from a stack is traditionally called:    1. delete C. pop    2. peek D. remove 2. Which of the following stack operations could result in stack underflow?    1. is\_empty C. push    2. pop D. Two or more of the above answers | | [  [ | C A | ]  ] |
| 1. Which of the following applications may use a stack?    1. A parentheses balancing program.    2. Keeping track of local variables at run time.    3. Syntax analyzer for a compiler.    4. All of the above. | | [ | D | ] |
| 28. Here is an infix expression: 4 + 3 \* (6 \* 3 - 12). Suppose that we are using the usual stack algorithm to convert the expression from infix to postfix notation. What is the maximum number of symbols that will appear on the stack AT ONE TIME during the conversion of this  expression? | | [ | D | ] |
| A. 1 | C. 3 | | | |
| B. 2 | D. 4 | | | |
| 1. What is the value of the postfix expression 6 3 2 4 + - \*    1. Something between -15 and -100    2. Something between -5 and -15    3. Something between 5 and -5    4. Something between 5 and 15    5. Something between 15 and 100 | | [ | A | ] |
| 30. If the expression ((2 + 3) \* 4 + 5 \* (6 + 7) \* 8) + 9 is evaluated with \*  having precedence over +, then the value obtained is same as the value of which of the following prefix expressions? | | [ | A | ] |
| A. + + \* + 2 3 4 \* \* 5 + 6 7 8 9 | C. \* + + + 2 3 4 \* \* 5 + 6 7 8 9 | | | |
| B. + \* + + 2 3 4 \* \* 5 + 6 7 8 9 | D. + \* + + 2 3 4 + + 5 \* 6 7 8 9 | | | |

1. Evaluate the following prefix expression:

+ \* 2 + / 14 2 5 1

[ B ]

* 1. 50
  2. 25
  3. 40
  4. 15

1. Parenthesis are never needed prefix or postfix expression: [ A ]
   1. True
   2. False
   3. Cannot be expected
   4. None of the above
2. A postfix expression is merely the reverse of the prefix expression: [ B ]
   1. True
   2. False
   3. Cannot be expected
   4. None of the above
3. Which among the following data structure may give overflow error, even though the current number of elements in it, is less than its size:

[ A ]

* 1. Simple Queue
  2. Circular Queue
  3. Stack
  4. None of the above

1. Which among the following types of expressions does not require precedence rules for evaluation:
   1. Fully parenthesized infix expression
   2. Prefix expression
   3. both A and B
   4. none of the above

[ C ]

1. Conversion of infix arithmetic expression to postfix expression uses: [ D ]
   1. Stack
   2. circular queue
   3. linked list
   4. Queue